# Counting Matrices with a Given Rank and Stable Rank over $M_{n}\left(\mathbb{F}_{p}\right)$ 

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- Crabb establishes a bijective correspondence between the set of nilpotent matrices in $M_{n}\left(\mathbb{F}_{p}\right)$ and the set of sequences of length $n-1$ of vectors of length $n$ over $\mathbb{F}_{p}$. He used this to deduces that the number of nilpotent matrices is $p^{n(n-1)}$.


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- We generalise Crabb's construction to a bijection involving all elements of $M_{n}\left(\mathbb{F}_{p}\right)$ and show how it can be used to count the numbers of matrices with specified stable rank and rank power sequence.


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## Examples

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A=\left(\begin{array}{llll}
0 & 1 & 3 & 6 \\
0 & 0 & 1 & 2 \\
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B^{3}=B^{4}=\ldots=\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0
\end{array}\right) \text { so B has stable rank 1 }
\end{gathered}
$$

## Crabb's Algorithm

Let M be an $n \times n$ nilpotent matrix of index $k$. Let $V_{0}, V_{1}, \ldots, V_{k}$ be ordered bases of the column spaces of $M^{0}, M^{1}, \ldots, M^{k}$. Let $V=V_{0}$ and "adapt" the basis V as follows. Starting with $\mathrm{i}=1$ :

- Let $W_{i}$ be the matrix that has the $j$ th vector of $V_{i}$ as its $j$ th row, over basis $V$. Reduce this to row echelon form.
- Let $c$ be the number of columns in $W_{i}$, and $p$ be the number of pivots. Let the first $c-p$ elements of the adapted basis $V^{\prime}$ be the $j$ th element of $V$, where $j$ takes the values of the columns without pivots. Then let the remaining elements be the sum of the $j$ th entry in a row multiplied by the $j$ th entry in $V$, for $j=1,2, \ldots, n$
- Then let $V=V^{\prime}$, and repeat the process for the next value of $i$, until $i=k$.

Now $V$ has been fully adapted. Finally, multiply each vector in $V$ on the left by $M$, and these are the entries in the tuple.

## Extending the Algorithm

We extended this algorithm to apply to all $n \times n$ matrices. In the extended algorithm, each $n \times n$ matrix has a corresponding $n$-tuple, instead of an ( $n-1$ )-tuple. We were then able to use this tuple to look at interesting properties of matrices, over a finite field of order $p$. For example, we were able to derive a formula to count the number of $n \times n$ matrices over a field of order $p$, with rank $r$ and stable rank $s$. In order to do this, we needed to define a few more terms:

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- The rank array of an $n \times n$ matrix $M$ is a non-decreasing sequence of integers. The $k$ th value of the rank array is the dimension of the span of the last $k$ entries in the tuple corresponding to $M$.


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- The rank array of an $n \times n$ matrix $M$ is a non-decreasing sequence of integers. The $k$ th value of the rank array is the dimension of the span of the last $k$ entries in the tuple corresponding to $M$.
- The power sequence of an $n \times n$ matrix $M$ is a strictly increasing sequence of integers $\left[\operatorname{rank}\left(M^{k}\right), \operatorname{rank}\left(M^{k-1}\right), \ldots, \operatorname{rank}\left(M^{2}\right), \operatorname{rank}(M)\right.$, $\left.\operatorname{rank}\left(I_{n}\right)\right]$, where $k$ is the index of $M$.


## Proposition

While examining this extended algorithm, we discovered the following:
An $n \times n$ matrix $M$ is nilpotent if the first entry in the rank array is zero. Otherwise, the stable rank is the first repeated entry in the rank array.

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- The number $M_{1}$ of $n \times n$ matrices over the finite field of order $p$ with rank $r$, and stable rank $s$ is as follows:

$$
M_{1}(\mathbf{n}, \mathbf{p}, \mathbf{r}, \mathbf{s})=\mathbf{p}^{(\mathbf{n}-r) s}\binom{n-s-1}{r-s} \prod_{p} \prod_{i=0}^{r-1}\left(p^{n}-p^{i}\right)
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- The number $M_{2}$ of $n \times n$ matrices over the finite field of order $p$ with rank $r$, stable rank $s$, and power sequence $L=\left[s=I_{1}, I_{2}, \ldots, I_{k-1}=\right.$ $\left.r, I_{k}=n\right]$ is as follows:

$$
M_{2}(n, p, r, s, L)=p^{s} \prod_{i=0}^{r-1}\left(p^{n}-p^{i}\right) \prod_{j=1}^{k-2} p^{l_{j+2}-2 I_{j+1}+I_{j}}\binom{l_{j+2}-I_{j+1}}{I_{j+1}-I_{j}}_{p}
$$

Crabb, M.C., Counting nilpotent endomorphisms, Finite Fields and Their Applications 12(151-154), 2005

回 Weisstein, Eric W., q-Binomial Coefficient, http://mathworld.wolfram.com/q-BinomialCoefficient.html

